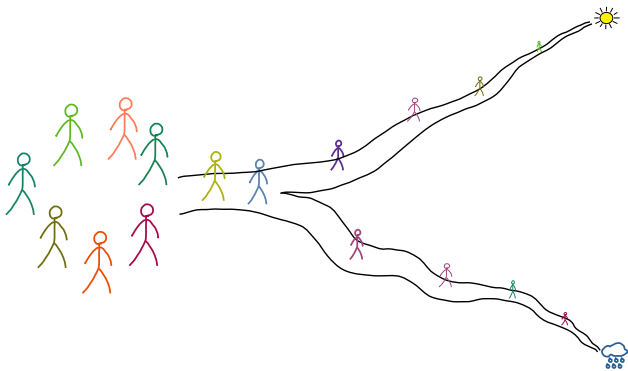


Public Projects with Preferences and Predictions¹



Bo Waggoner
University of Colorado, Boulder

ESSET, Gerzensee
July 2024

¹Based on joint work with Mary Monroe. Supported by the Ethereum Foundation.



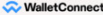
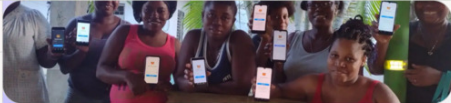




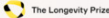



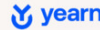



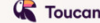
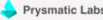



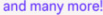

Inspiration: blockchain governance

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Gitcoin (2017-):

- \$60M distributed
- Uses “quadratic funding” donation-matching

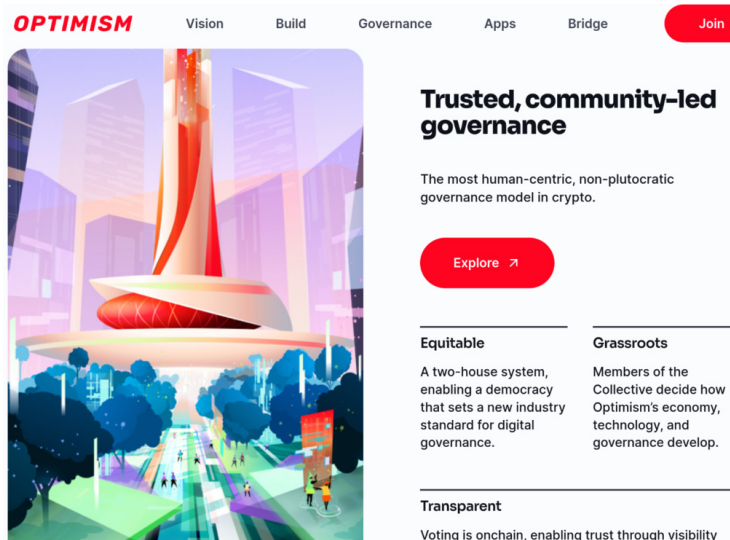
RAISING ON GITCOIN

Inspiration: blockchain governance

Layer 2's and platforms:

- e.g. Uniswap (2018-), Arbitrum (2018/2021-), Optimism (2019-)



OPTIMISM Vision Build Governance Apps Bridge [Join](#)

Trusted, community-led governance

The most human-centric, non-plutocratic governance model in crypto.

[Explore ↗](#)

Equitable
A two-house system, enabling a democracy that sets a new industry standard for digital governance.

Grassroots
Members of the Collective decide how Optimism's economy, technology, and governance develop.

Transparent
Voting is onchain, enabling trust through visibility

Inspiration: blockchain governance

Layer 2's and platforms:

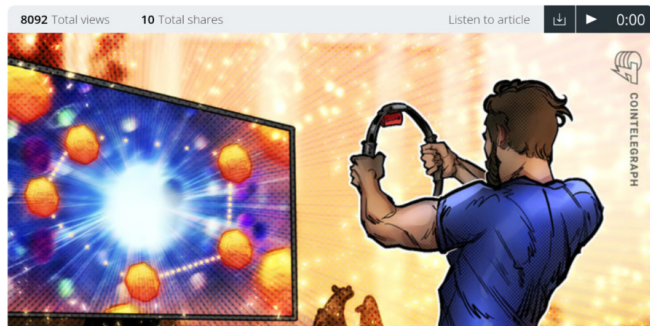
- e.g. Uniswap (2018-), Arbitrum (2018/2021-), Optimism (2019-)

AMAKA NWAOKOCHA

JUN 09, 2024

Arbitrum to distribute \$215M in ARB tokens for gaming innovation

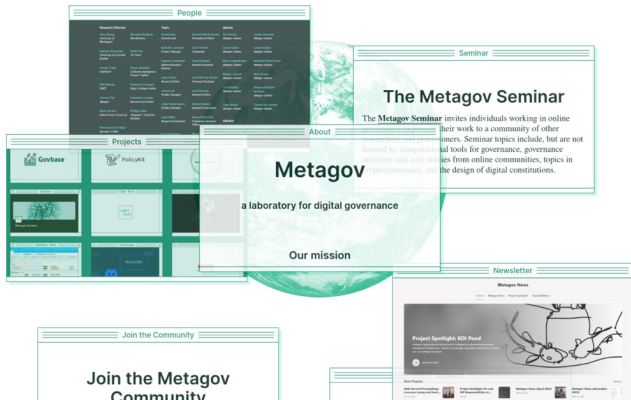
Initially introduced in March, the proposal was approved on June 7, with over 75% of votes in favor.



Inspiration: blockchain governance

Non-mechanism design governance research, e.g. at CU Boulder:

- Nathan Schneider: co-ops perspective, e.g. Metagov
- Eric Alston: government & corporation perspective, e.g. constitutions



Paradigms for group decisionmaking

Paradigms for group decisionmaking

aggregate
preferences

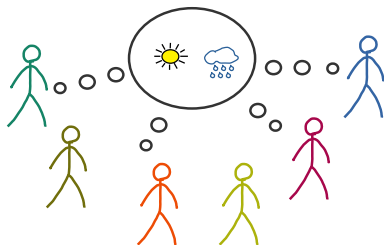


Paradigms for group decisionmaking

aggregate preferences



aggregate information

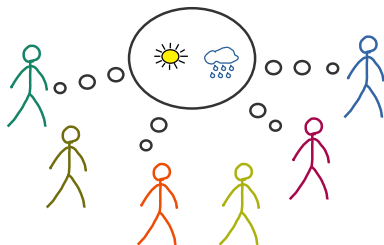


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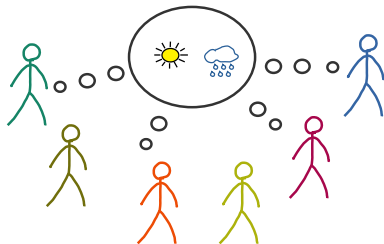
Can a formal mechanism **do both**?

Paradigms for group decisionmaking

aggregate preferences



aggregate information



Can a formal mechanism **do both**?

Hanson ("futarchy", 2000; 2007); Schoenebeck and Tao (2021);
Amanatidis, Birmpas, Lazos, and Marmolejo-Cossío (2022)

Setting: public projects with organizational mission

voters



mission

e.g. minimize
carbon footprint

Setting: public projects with organizational mission

voters

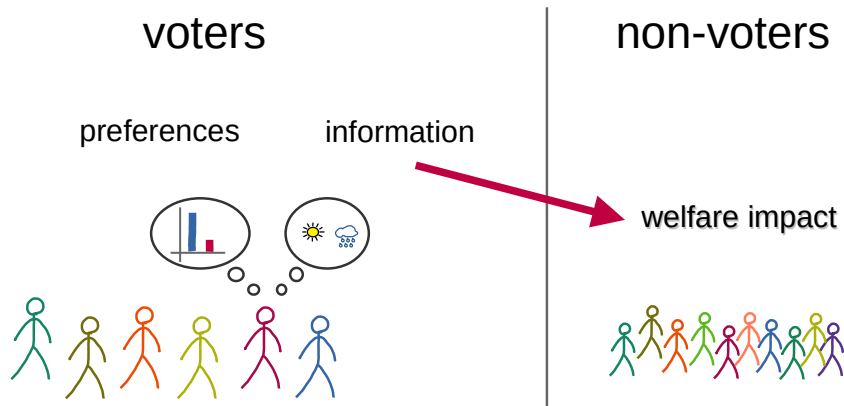


non-voters

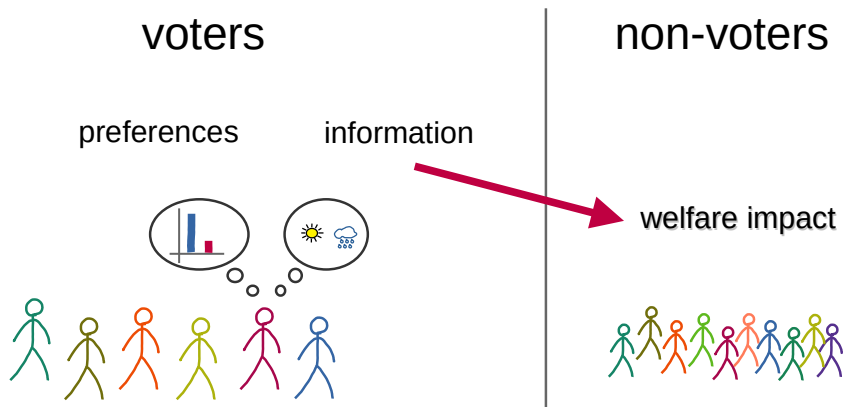
welfare impact



Setting: public projects with organizational mission



Setting: public projects with organizational mission



Goal: welfare guarantees ("Price of Anarchy")

Outline:

- 1 Public Projects from preferences
- 2 Public Projects from predictions
- 3 Public Projects with preferences and predictions

Outline:

1 Public Projects from preferences

- Model, definitions
- Related work: VCG, QTM
- Our results: QTM

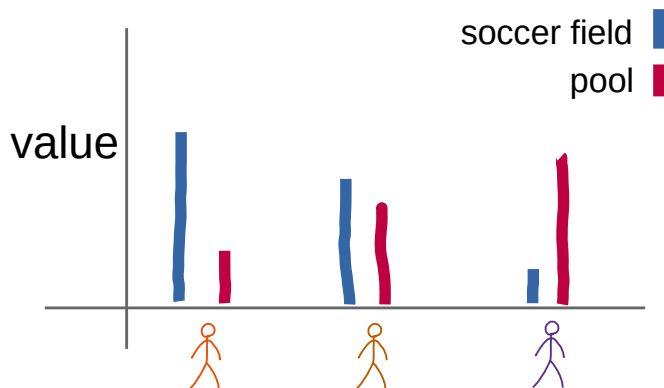
2 Public Projects from predictions

3 Public Projects with preferences and predictions

Public projects: model

Each agent i has a value v_k^i for each option k

nonnegative

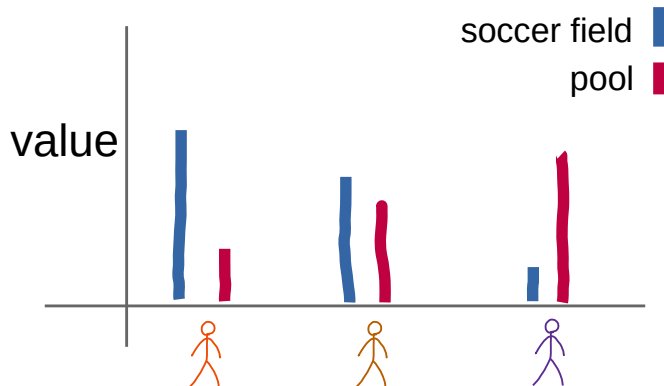


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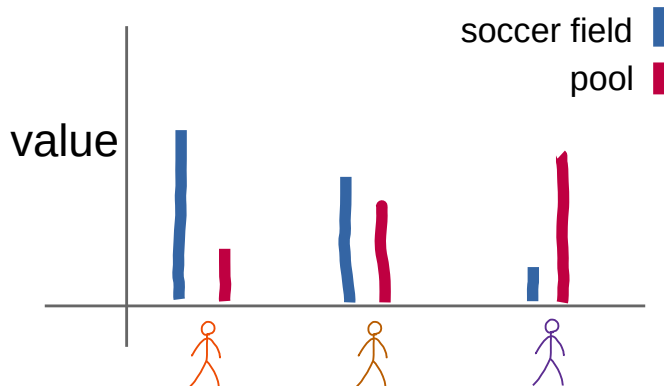


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- Mechanism: collect reports, pick an alternative k , assign payments



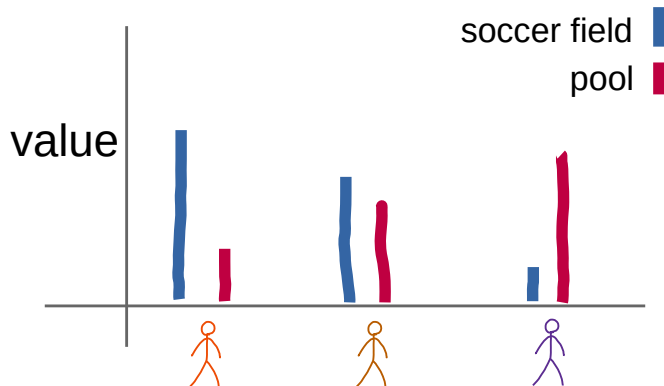
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quasilinear



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Price of Anarchy

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quasilinear

$$\frac{\mathbb{E}[V_k]}{\max_{k'} V_{k'}}$$

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quasilinear

$$\text{Price of Anarchy} = \min_{\text{equilibria}} \frac{\mathbb{E}[V_k]}{\max_{k'} V_{k'}}$$

*our mechanisms: pure-strategy Nash equilibria
(convex strategy space, strictly concave utilities)*

Related work

- **VCG mechanism:** Price of Anarchy = 0
Not budget-balanced, revenue unstable

efficient equil. exists

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- **“First-price”**:² Price of Anarchy $\rightarrow 0$ *but sequential model: = 1*

²Lucier, Singer, Syrgkanis, Tardos (2013)

Related work

- **VCG mechanism:** Price of Anarchy = 0 *efficient equil. exists*
Not budget-balanced, revenue unstable
- **“First-price”:**² Price of Anarchy $\rightarrow 0$ *but sequential model: = 1*
- **Quadratic Transfers Mechanism:**³
In an i.i.d. model, Price of Anarchy $\rightarrow 1$
as population grows large

²Lucier, Singer, Syrgkanis, Tardos (2013)

³Eguia, Immorlica, Ligett, Weyl, Xefteris (2019; 2023).

Quadratic Transfers Mechanism (QTM)⁴

Agent i submits votes $\{a_k^i\}$

⁴Eguia, Immorlica, Ligett, Weyl, Xefteris (2019; 2023).

Quadratic Transfers Mechanism (QTM)⁴

Agent i submits votes $\{a_k^i\}$ and pays $c \sum_k (a_k^i)^2$
can redistribute payment to all others \implies budget-balanced

$c = a$ parameter

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Pick $k = \arg \max A_k$ (?)

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~~Pick $k = \arg \max A_k$ (?)~~

Pick $k \sim \mathbf{p}$ randomly where

“soft max”

$$p_k = \frac{e^{A_k}}{e^{A_1} + \dots + e^{A_m}}.$$

⁴Eguia, Immorlica, Ligett, Weyl, Xefteris (2019; 2023).

Our results on QTM for public projects

Theorem (Monroe and Waggoner (2024))

Let $v^* = \max_{i,k} v_k^i$ and $\epsilon = \frac{v^*}{\max_k V_k}$.

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For two alternatives, choosing $c = \frac{1}{2}v^*$, the QTM has

$$\text{Price of Anarchy} \geq \max \left\{ \frac{1}{2}, 1 - (2\epsilon)^{2/5} \right\}.$$

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Notes:

- builds on tools of analogous asymptotic result of Eguia et. al
- 3+ alternatives: $\text{PoA} \geq \frac{1}{\# \text{ alternatives}}$; better is open nonasymptotically.

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- 2 **Public Projects from predictions**
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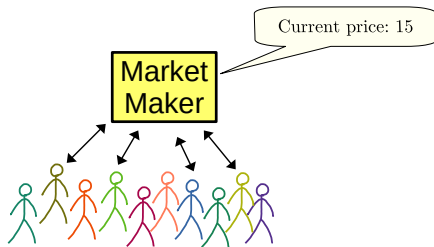
Outline:

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- 2 **Public Projects from predictions**
 - Prediction markets
 - Decision markets
- 3 Public Projects with preferences and predictions

Background: prediction markets

Prediction markets:

all we need today



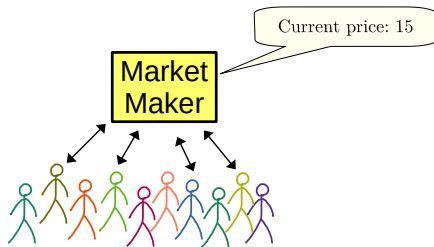
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- Want to predict X , e.g. tons of carbon emitted by city in 2025
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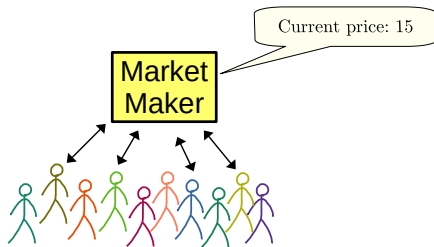
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 \implies get consensus prediction (price) \hat{X}
- On Jan 1, 2026: observe X , settle bets

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Ostrovsky (2012): In any equilibrium, all* information is aggregated:

$$\hat{X} = \mathbb{E}[X \mid \text{signals}]$$

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Alternative: wagering mechanisms

can average the predictions, but aggregation is not guaranteed

Decision markets⁵

Mechanism (pick among m alternatives):

1 Suppose $B_k =$ welfare impact of k

nonnegative, higher is better

⁵Hanson (1999); Othman and Sandholm (2010)

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Chen et al. (2010): **randomization** + importance weighting \implies truthful

Combine with Ostrovsky (2012): **approximately** efficient

⁵Hanson (1999); Othman and Sandholm (2010)

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- 2 Public Projects from predictions
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 - Model
 - Mechanism: SQUAP
 - Results
 - Caveats

voters



non-voters

total welfare if k : B_k



voters

private
values $\{v_k^i\}$

private signal
about $\{B_k\}$



non-voters

total welfare if k : B_k



voters

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non-voters

total welfare if k : B_k



$$W_k := V_k + \mathbb{E}[B_k \mid \text{signals}]$$

total welfare of option k

Related work

VCG+scoring rules mechanism of Cai, Mahdian, Mehta, Waggoner (2013)

Related work

VCG+scoring rules mechanism of Cai, Mahdian, Mehta, Waggoner (2013)

- Each i submits valuation $\{v_k^i\}$ and conditional predictions $\{p_k^i\}$
- Compute $\hat{B}_k = g(\{p_k^i\})$ *assume g component-wise convex*
- Select $k = \arg \max_k (V_k + \hat{B}_k)$
- Use VCG payments combined with scoring rules *constructed from g*

Exists **fully efficient** equilibrium

assuming you know how to aggregate

but **PoA** = 0

inherits VCG weaknesses

Mechanism

Synthetic-Players Quadratic Transfer Mechanism with Predictions (SQUAP):

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- Run QTM, but add “synthetic player” with values $(\hat{B}_1, \dots, \hat{B}_m)$
results in $k \sim p$
- Cancel all conditional markets but k
- Later, observe B_k and pay out k market
use importance-weighted payment of Chen et. al (2011)

Main result

Theorem (Monroe and Waggoner (2024))

Let $v^* = \max_{i,k} v_k^i$ and $\epsilon = \frac{v^*}{\max_k W_k}$.

$\epsilon =$ “influence”

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Assume “markets aggregate information” (A1).

For two alternatives, there is a choice of SQUAP parameters s.t.

$$\text{Price of Anarchy} \geq 1 - 2\epsilon - (2\epsilon)^{2/5}.$$

Assumption (A1):

Main result

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Assumption (A1):

market converges to $\mathbb{E}[B_k \mid \text{signals}]$, then manipulation occurs

or: markets aggregate information off the equilibrium path

or: nobody has **exclusive** private information

Main result

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Key intuitions (manipulation doesn't hurt much):

- Manipulating predictions is more costly than manipulating votes
- Importance weights: manipulation does not improve market payouts

Main result

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Notes:

- Can use revenue of QTM to subsidize prediction market, sometimes
result: QTM revenue = Θ (“disagreement”)
- Can use wagering instead of prediction markets
strategically easier, but assume aggregation is possible

The giant caveat

Unfortunately: you can't run SQUAP.

synthetic player needs knowledge of values to find equilibrium

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Unfortunately: you can't run SQUAP.

synthetic player needs knowledge of values to find equilibrium

- For that matter, can agents play QTM?
just need to respond to vote totals $\{A_k\}$, mean-field style
- Possible solution: run process over time with aim of convergence

Proposed variant: given \hat{B}_1, \hat{B}_2 , collect votes and pick using

$$p_1 = \frac{e^{A_1 + \frac{p_1 p_2}{v^*} (\hat{B}_1 - \hat{B}_2)}}{e^{A_1 + \frac{p_1 p_2}{v^*} (\hat{B}_1 - \hat{B}_2)} + e^{A_2 + \frac{p_1 p_2}{v^*} (\hat{B}_1 - \hat{B}_2)}}.$$

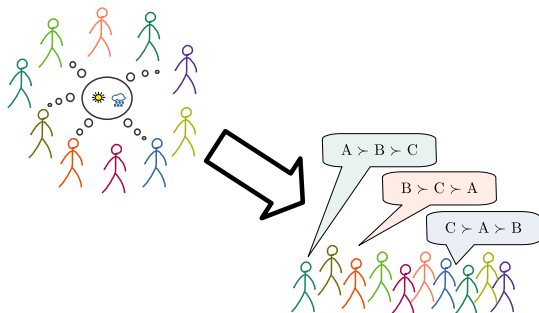
difficult to analyze, involves fixed-point computation

Future work: voters learn their preferences

In our model: Voters had *fixed* preferences.

Ideally: voters adjust preferences in response to aggregated information.⁶

Issue: market manipulation \implies misled voters \implies changed outcome.



⁶See Schoenebeck and Tao (2021)

Summary:

- Decisions should aggregate both preferences and information
- Proposed SQUAP, combining prediction markets and quadratic voting
- Proved Price of Anarchy bounds (under impractical assumptions)

Open:

- Analysis of “practical SQUAP”
- Better synthesis of information and preference aggregation
- Role of such mechanisms in a governance structure
- Can organizations avoid capture?

End

Summary:

- Decisions should aggregate both preferences and information
- Proposed SQUAP, combining prediction markets and quadratic voting
- Proved Price of Anarchy bounds (under impractical assumptions)

Open:

- Analysis of “practical SQUAP”
- Better synthesis of information and preference aggregation
- Role of such mechanisms in a governance structure
- Can organizations avoid capture?

Thanks!